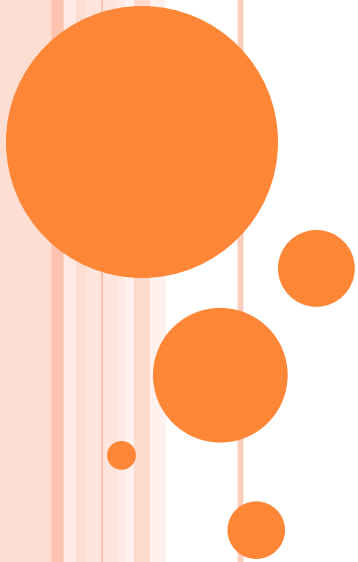
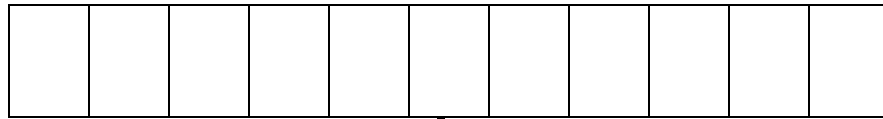


PUSHDOWN AUTOMATA

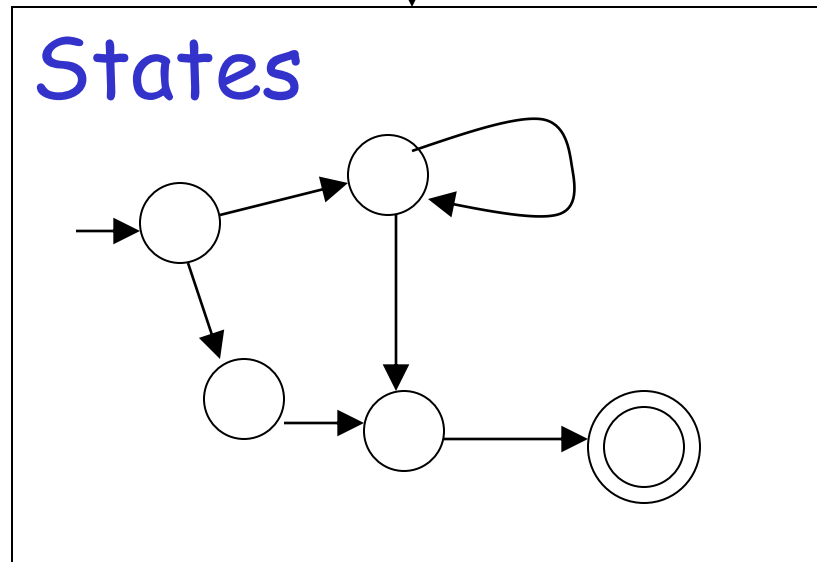
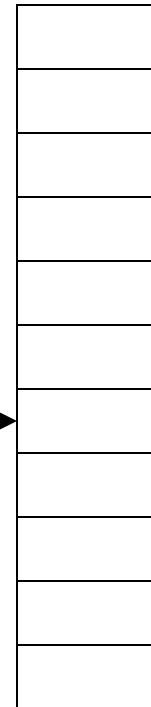


PUSHDOWN AUTOMATON -- PDA

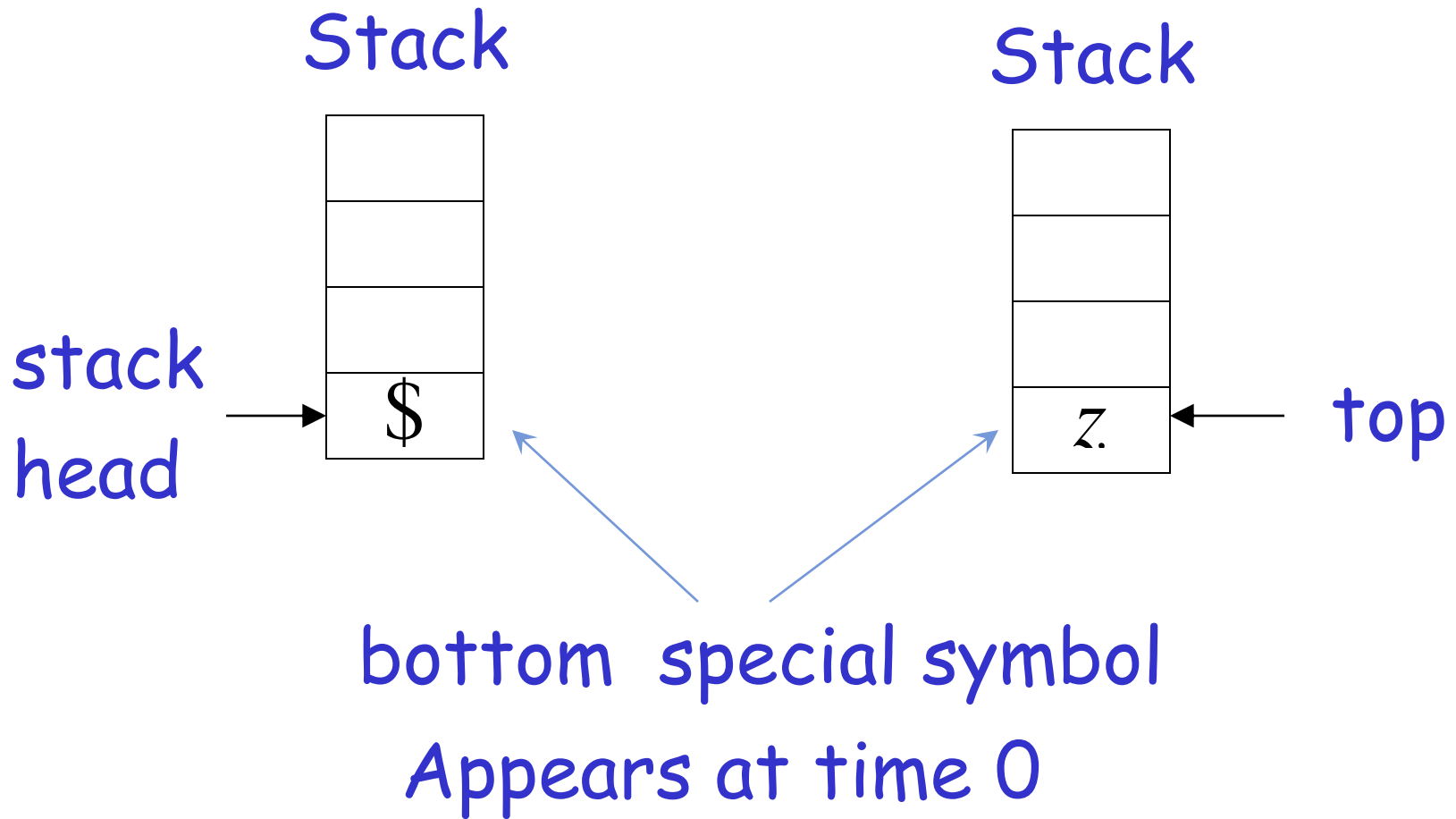
Input String



Stack



Initial Stack Symbol

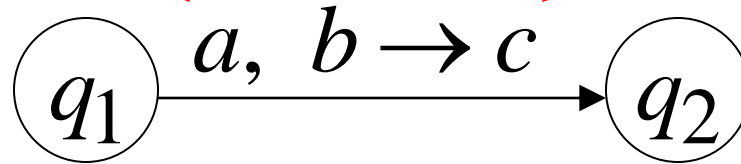


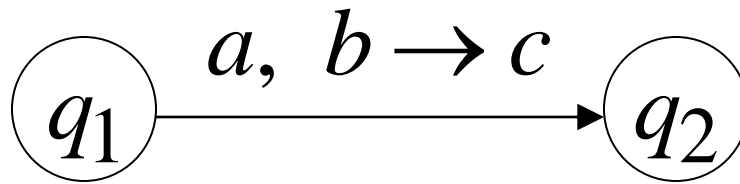
THE STATES

Input
symbol

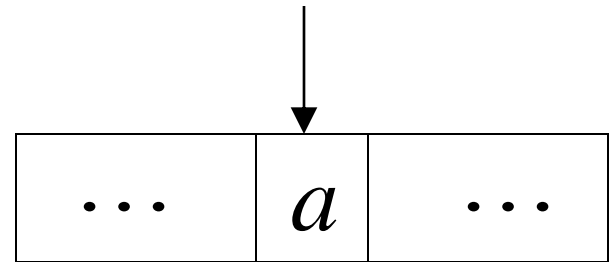
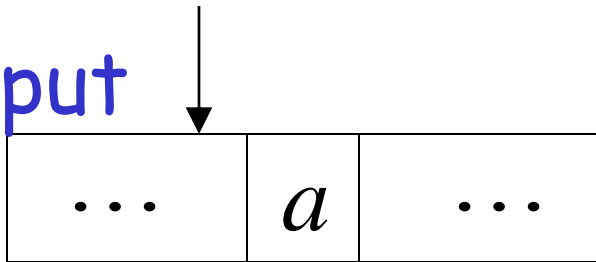
Pop
symbol

Push
symbol

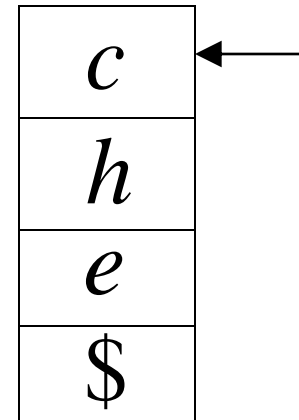
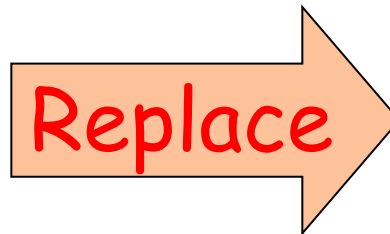
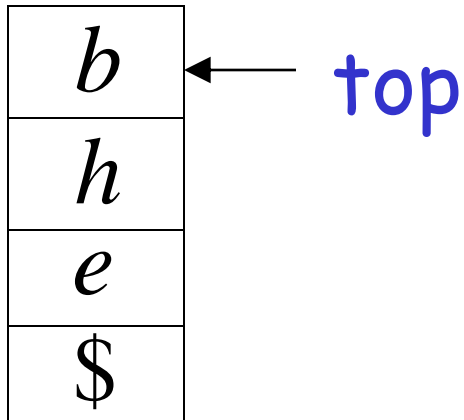


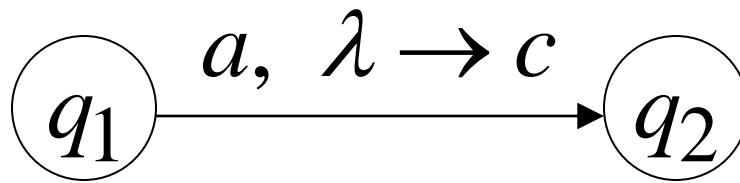


input

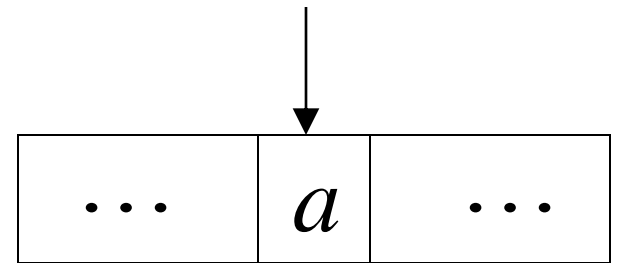
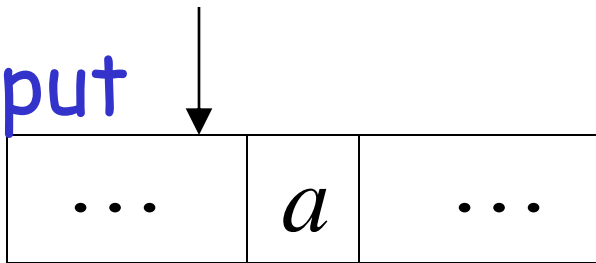


stack

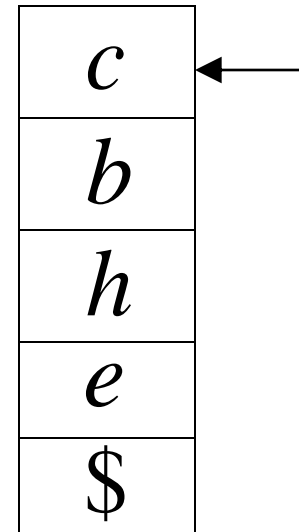
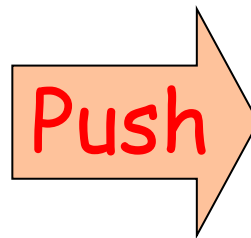
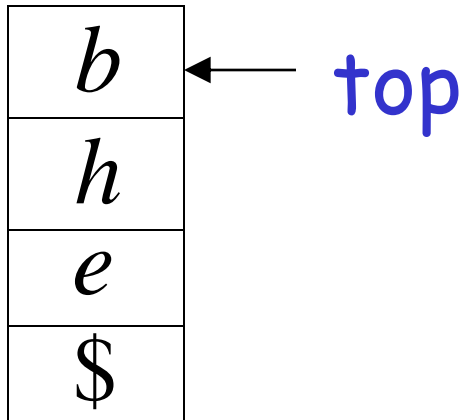


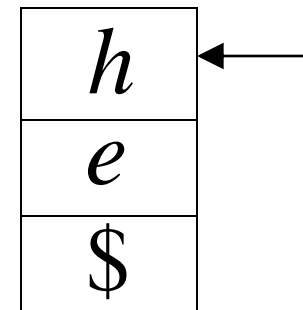
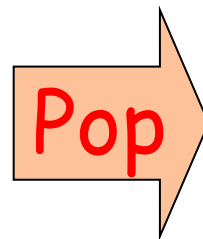
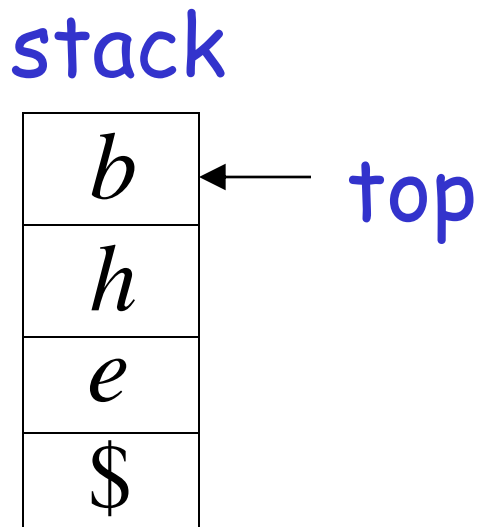
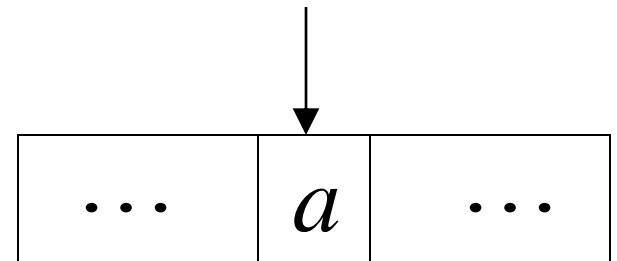
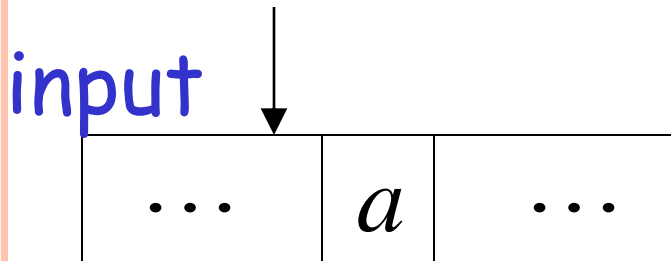
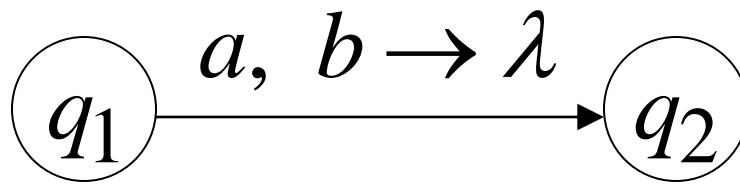


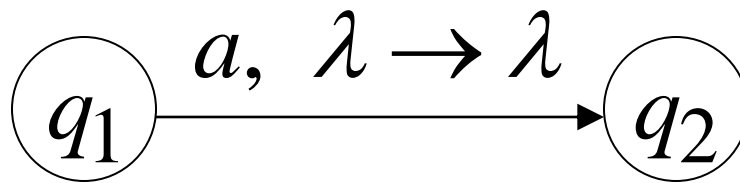
input



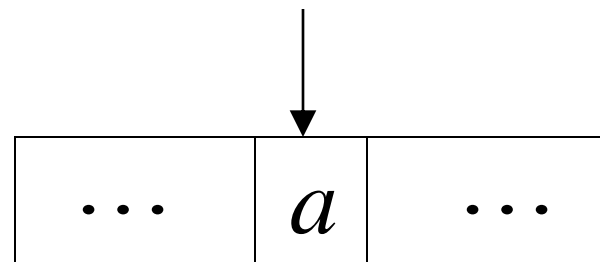
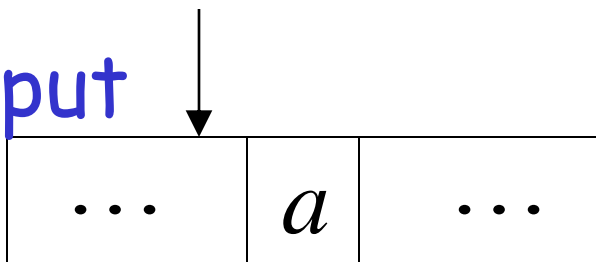
stack



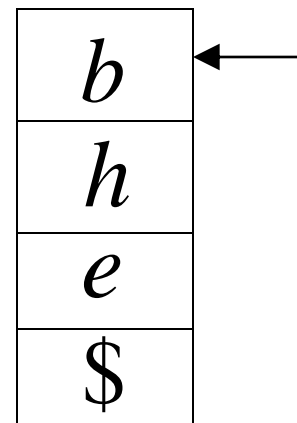
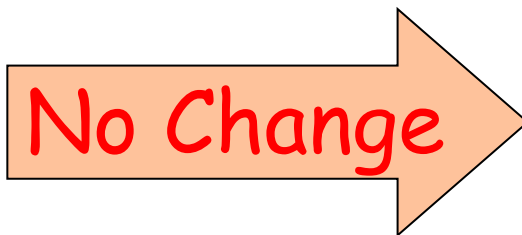
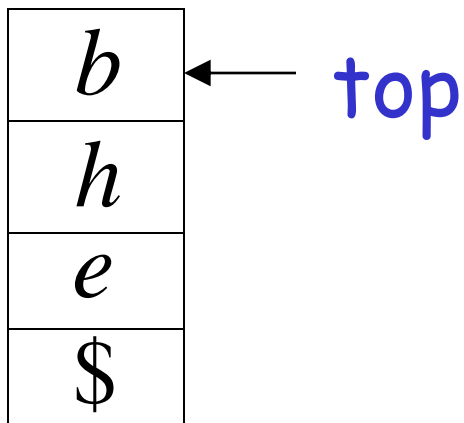




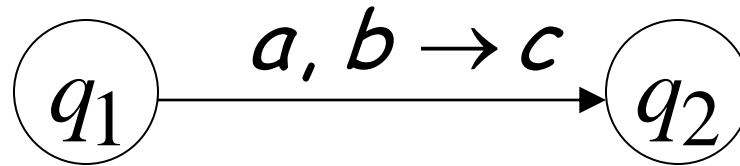
input



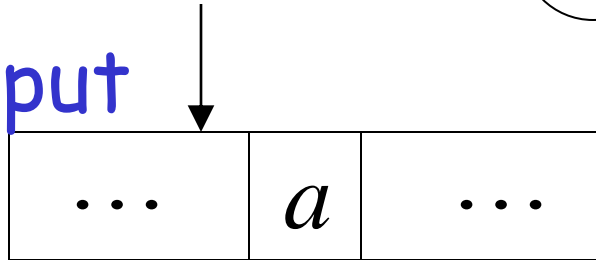
stack



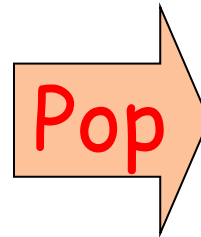
Pop from Empty Stack



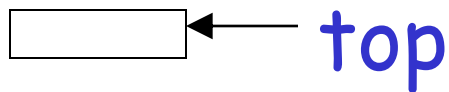
input



stack



Automaton halts!

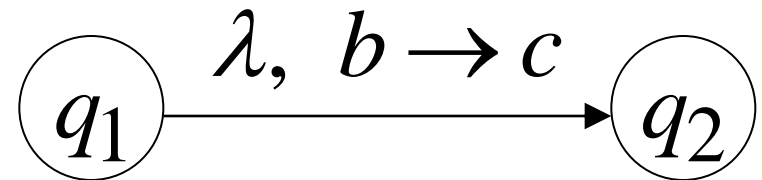
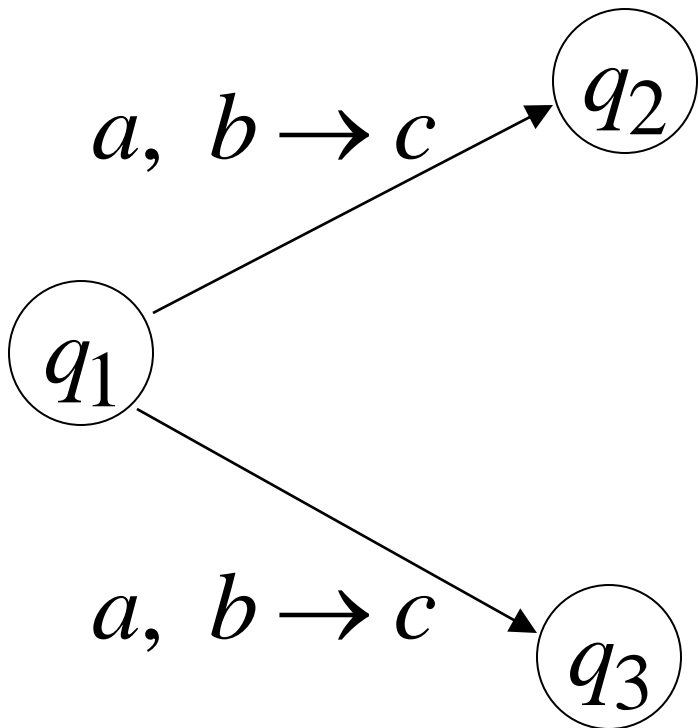


If the automaton attempts to pop from empty stack then it halts and rejects input

NON-DETERMINISM

PDAs are non-deterministic

Allowed non-deterministic transitions

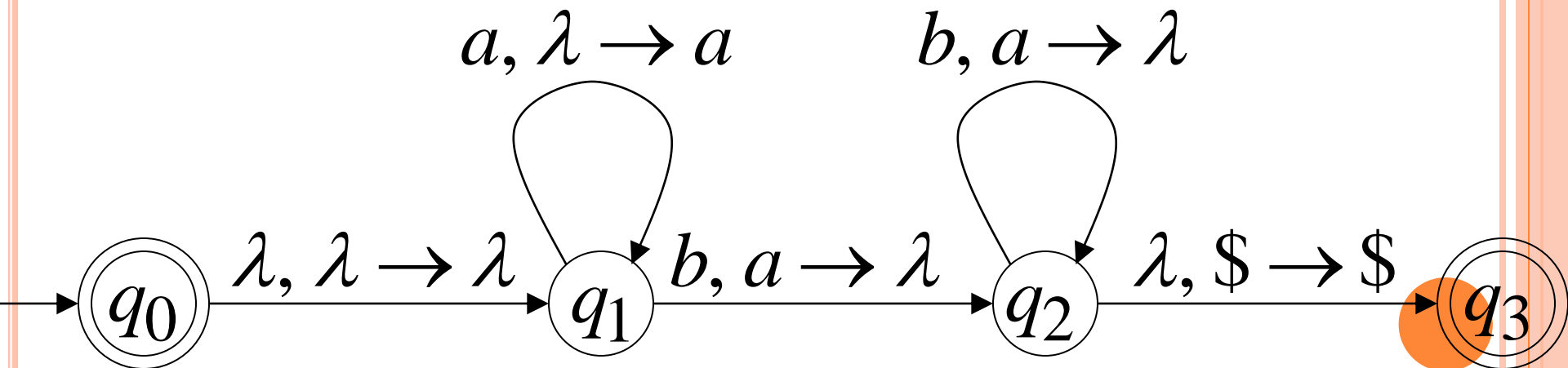


λ - transition



EXAMPLE PDA

PDA M : $L(M) = \{a^n b^n : n \geq 0\}$



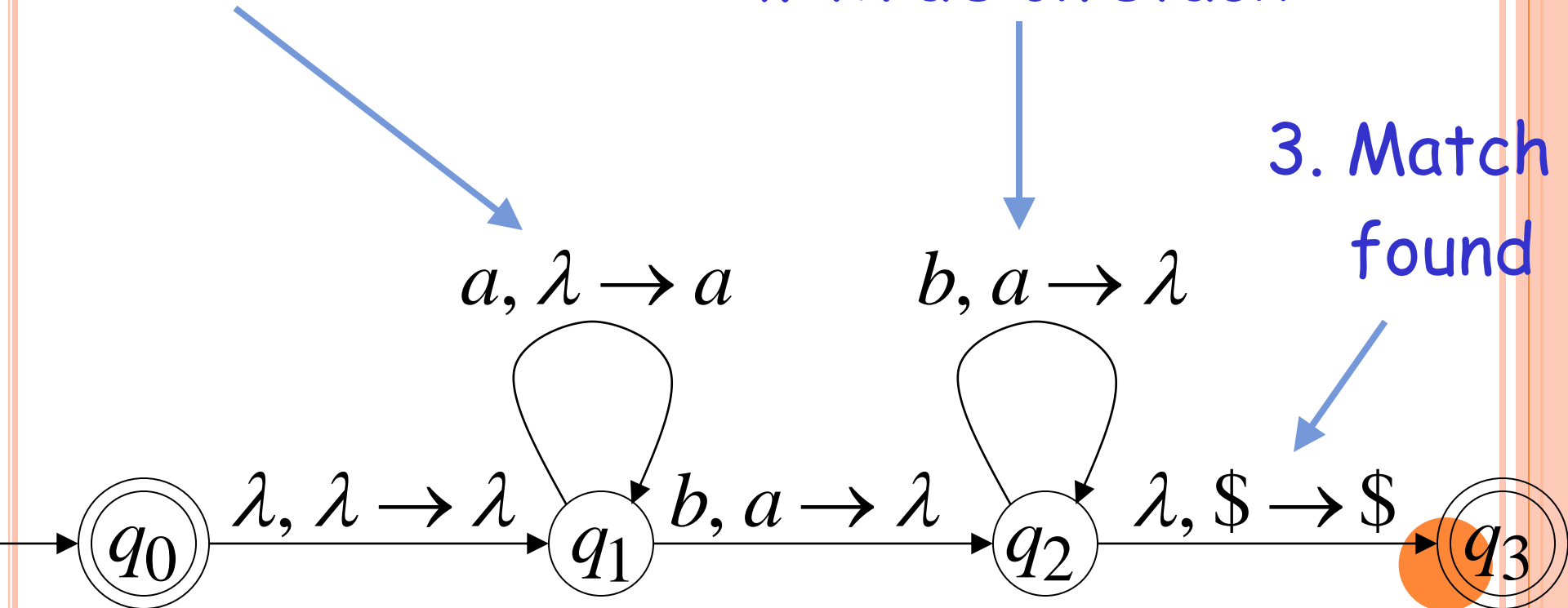
$$L(M) = \{a^n b^n : n \geq 0\}$$

Basic Idea:

1. Push the a's on the stack

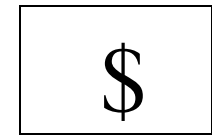
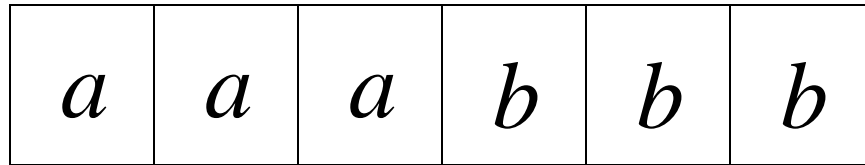
2. Match the b's on input with a's on stack

3. Match found



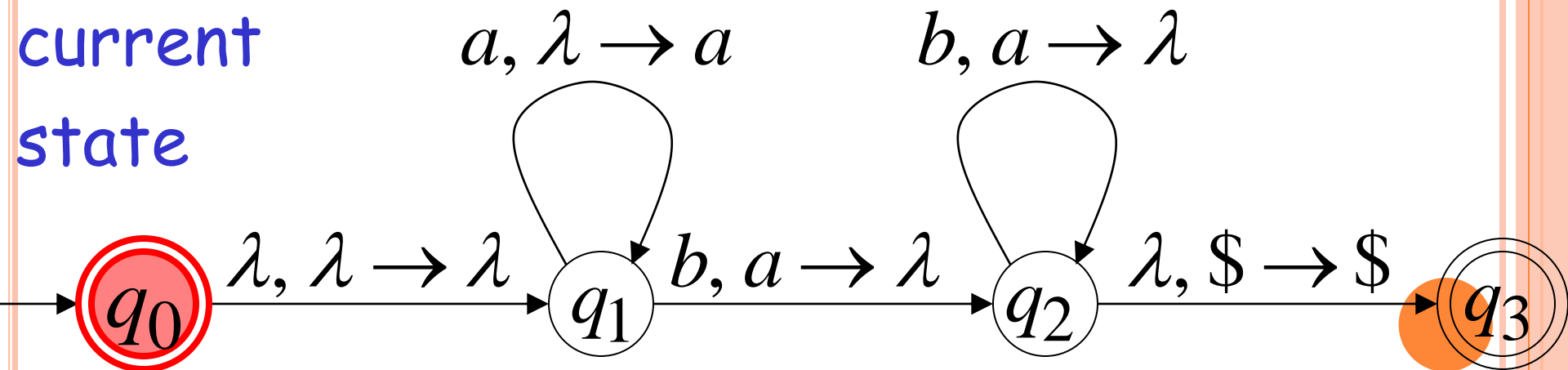
Execution Example: Time 0

Input



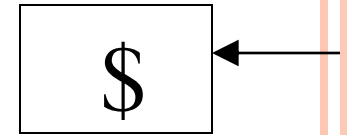
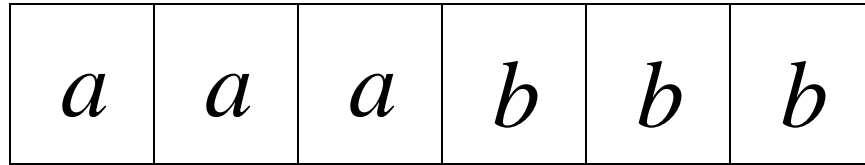
Stack

current
state

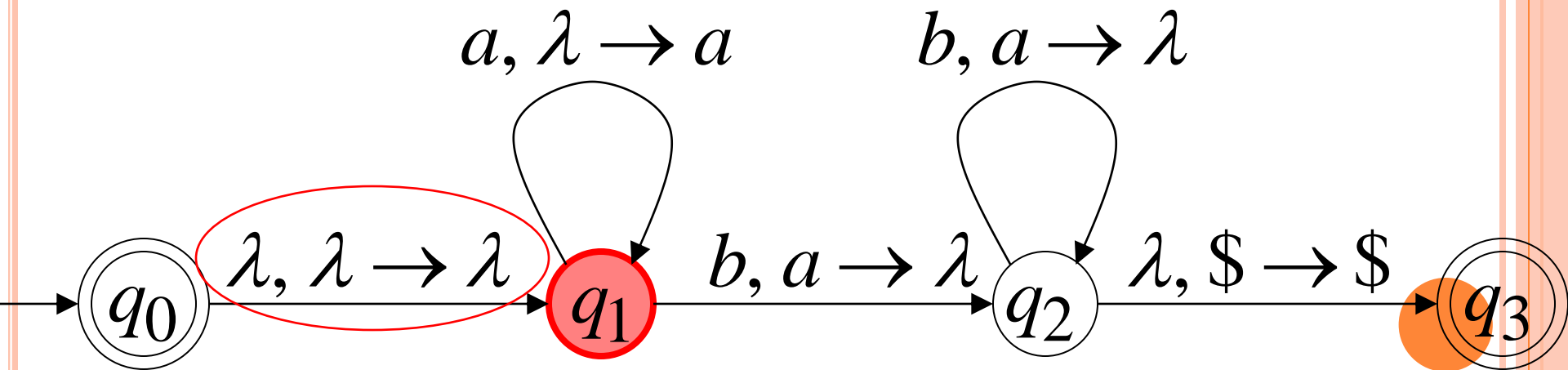


Time 1

Input

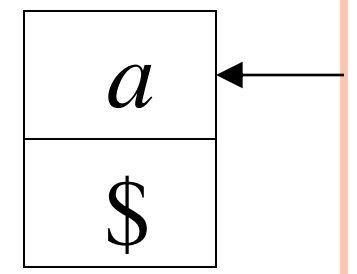
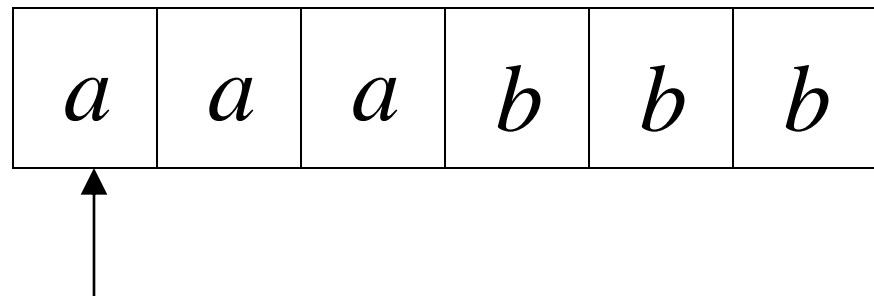


Stack

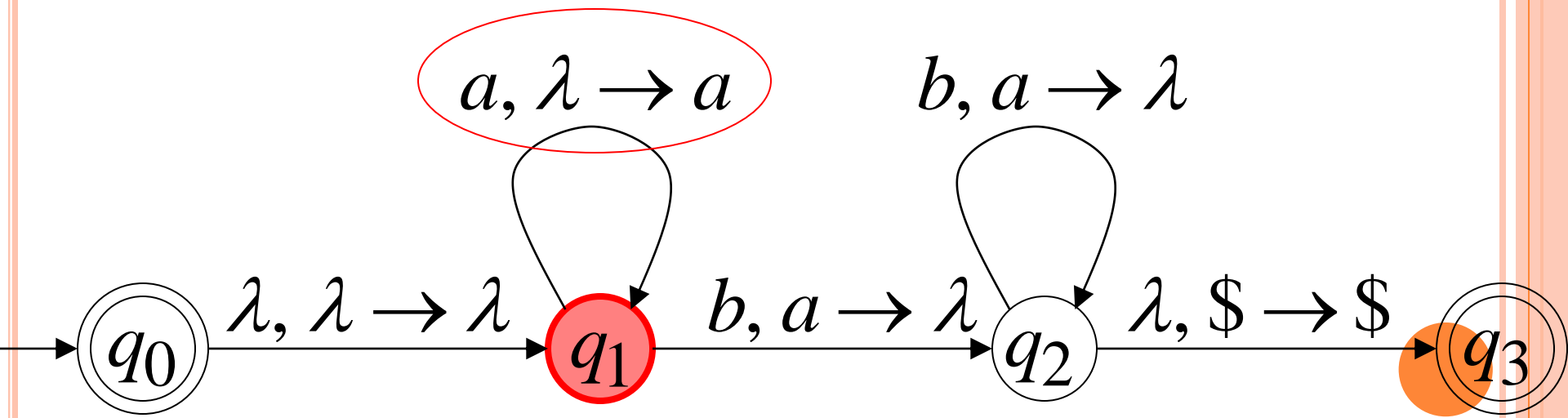


Time 2

Input

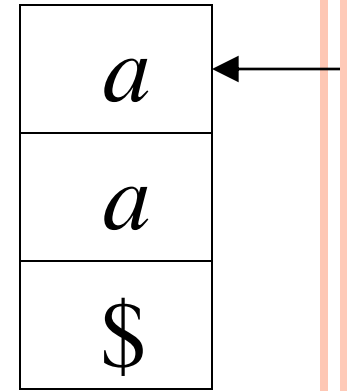
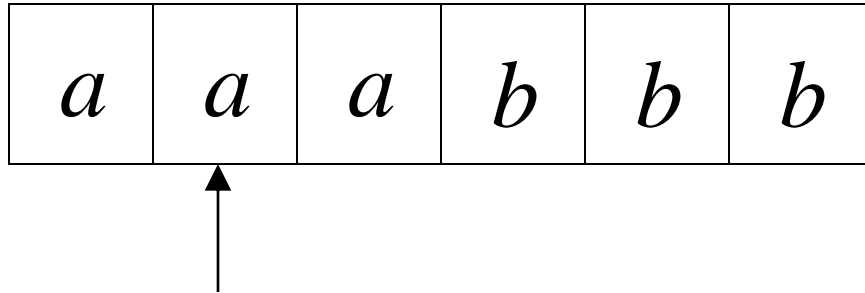


Stack

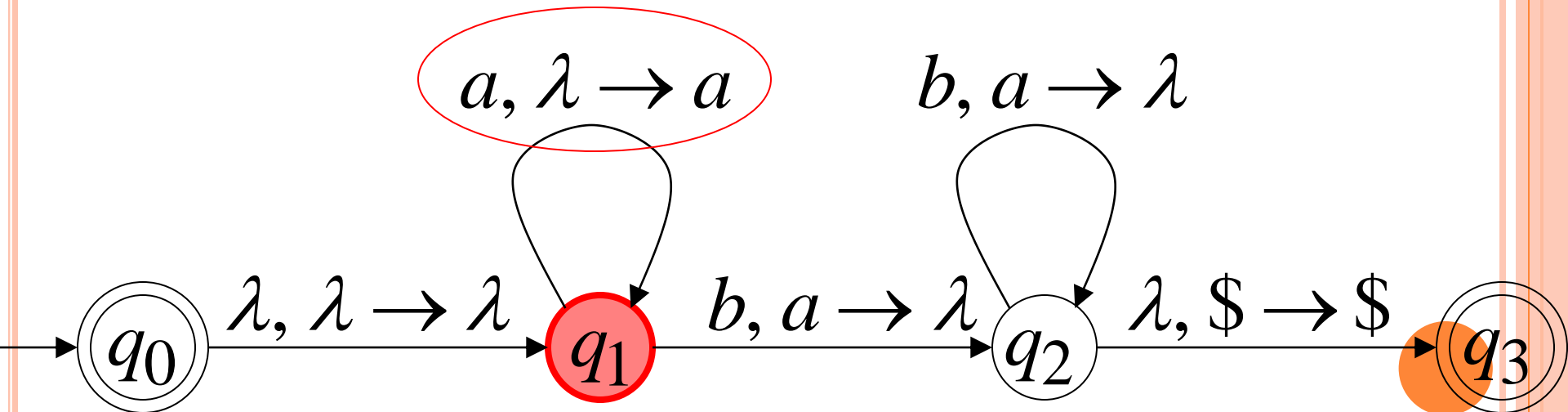


Time 3

Input

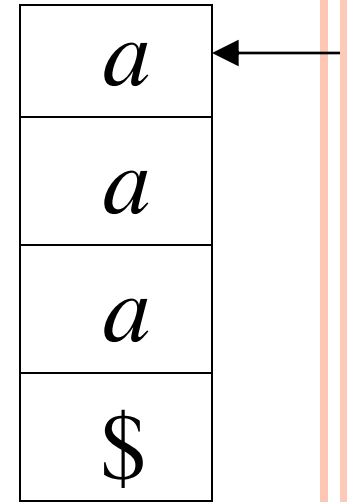
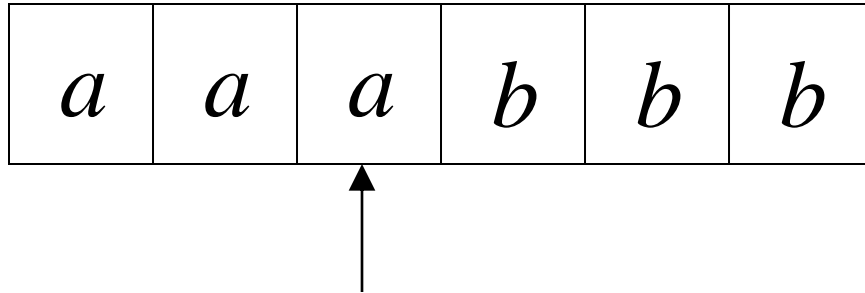


Stack

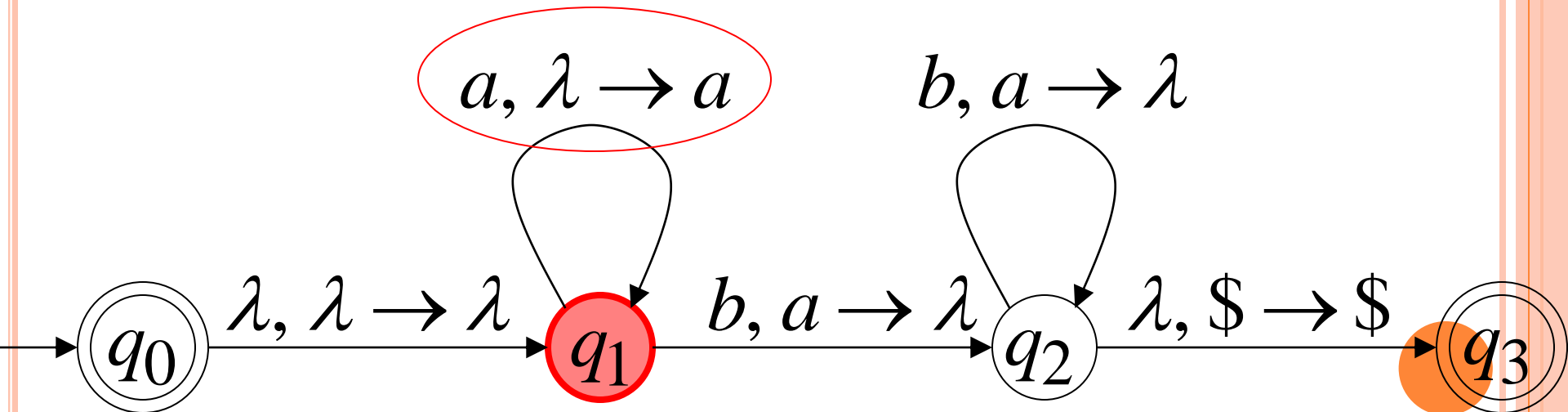


Time 4

Input

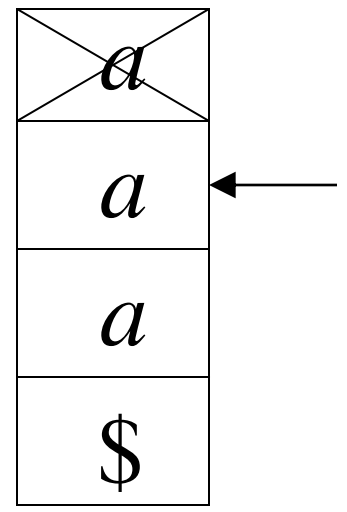
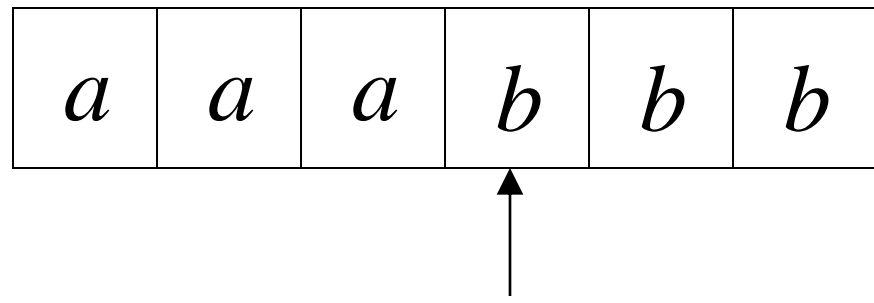


Stack

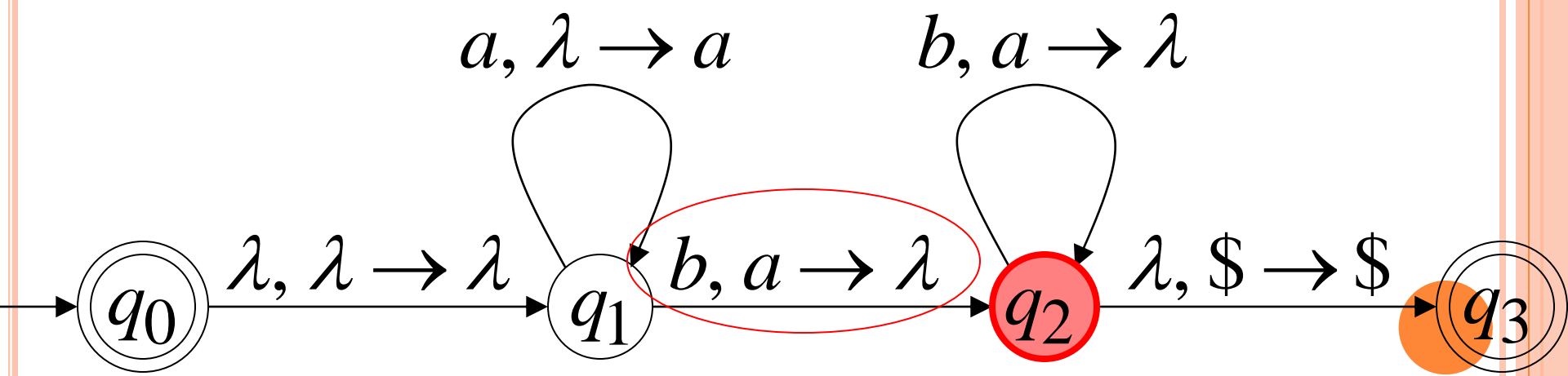


Time 5

Input

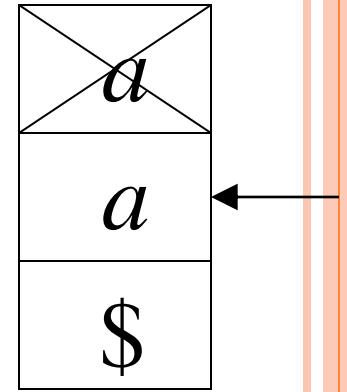
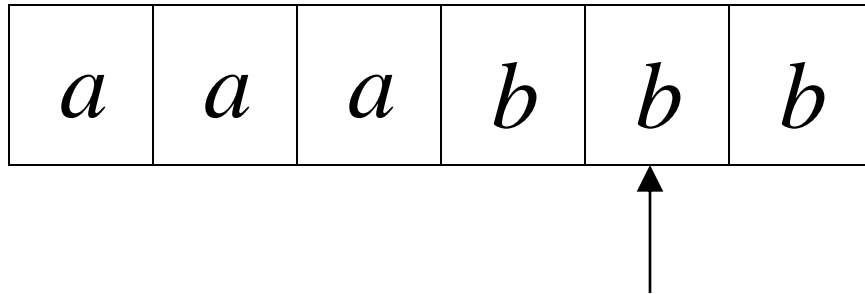


Stack

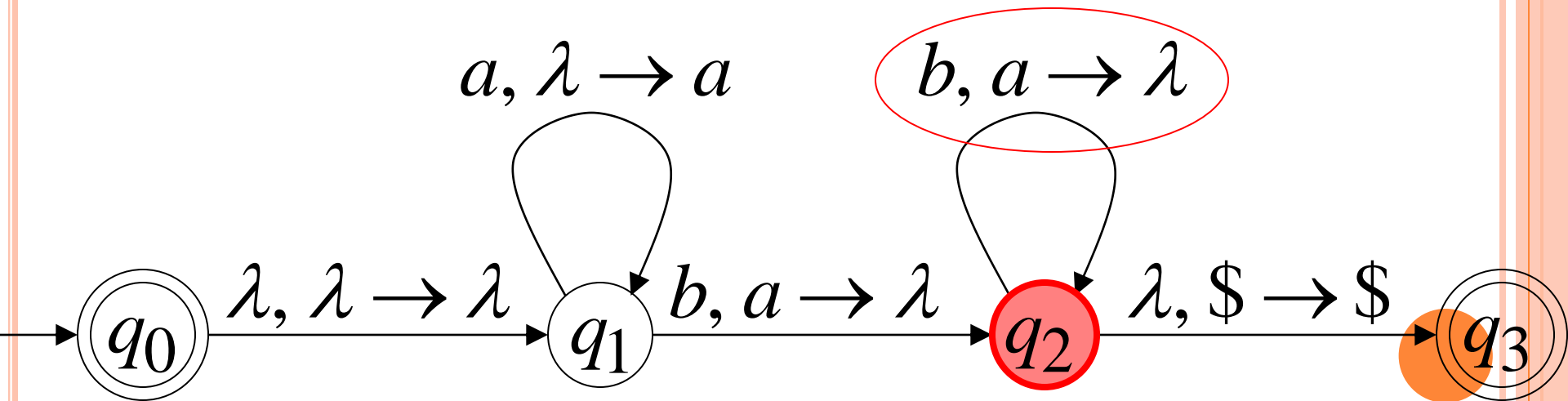


Time 6

Input

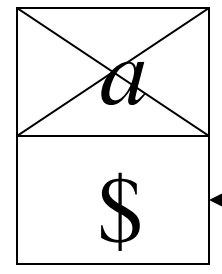
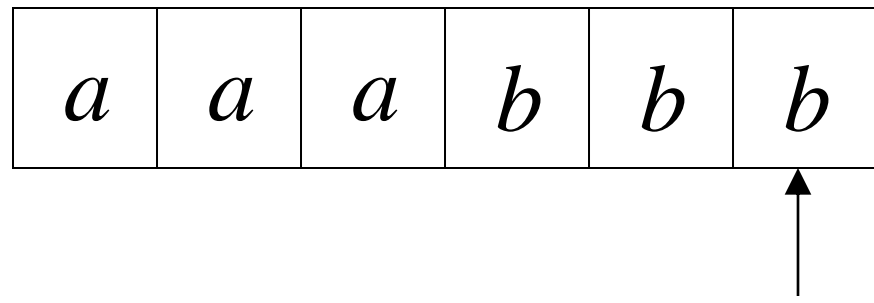


Stack

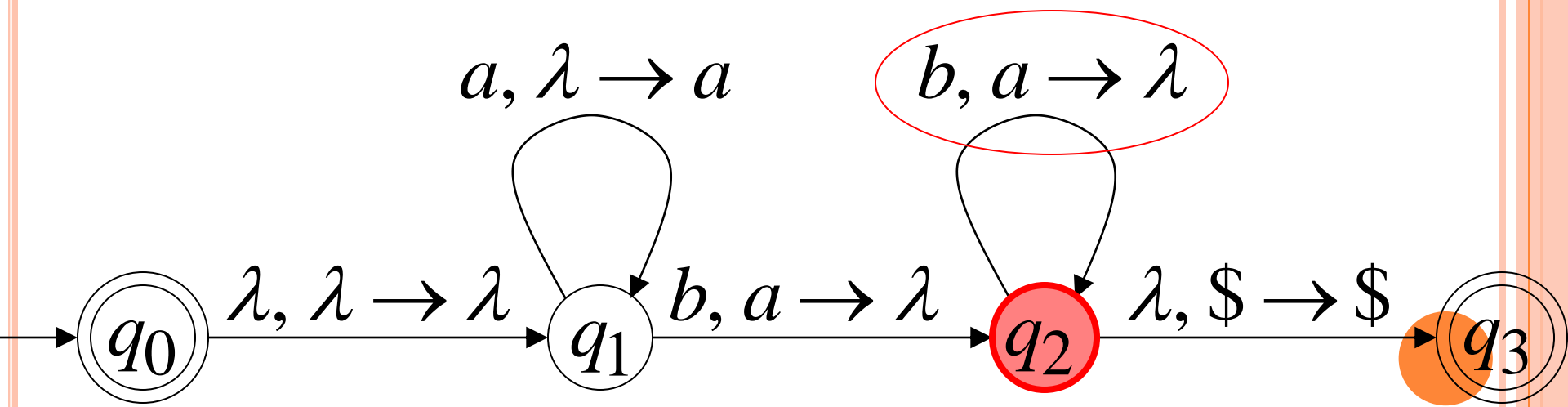


Time 7

Input

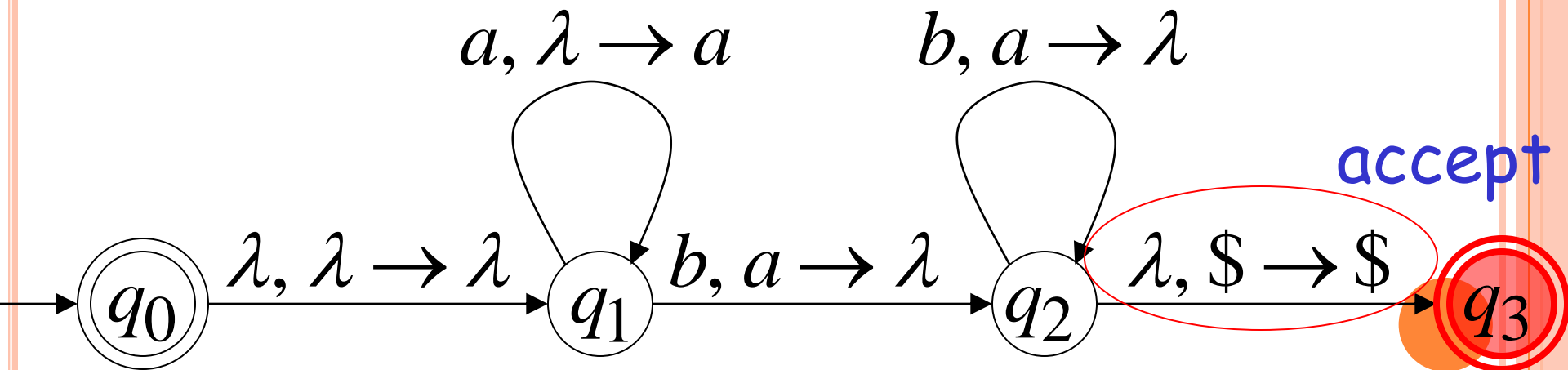
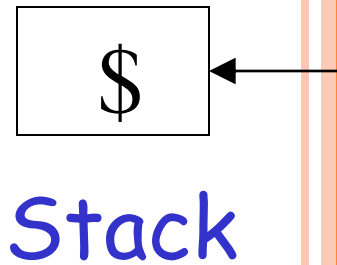
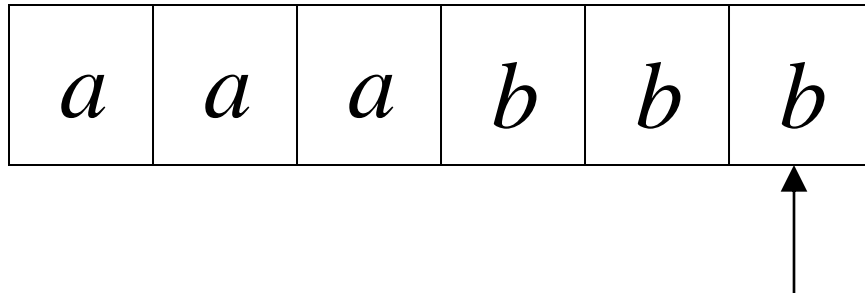


Stack



Time 8

Input



A string is accepted if there is
a computation such that:

All the input is consumed

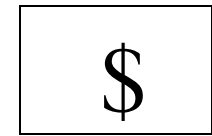
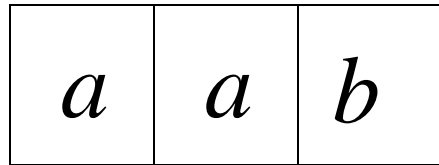
AND

The last state is an accepting state

we do not care about the stack contents
at the end of the accepting computation

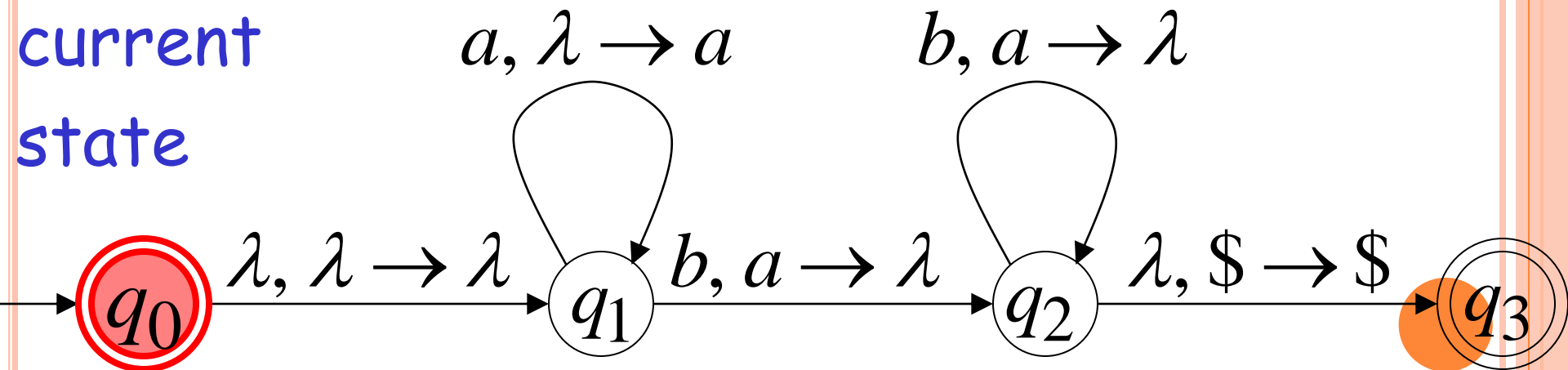
Rejection Example: Time 0

Input



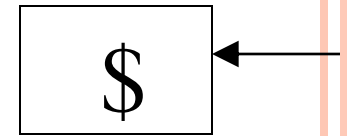
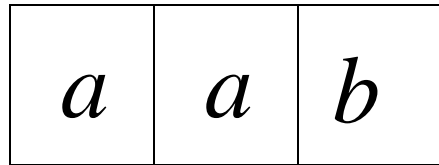
Stack

current
state



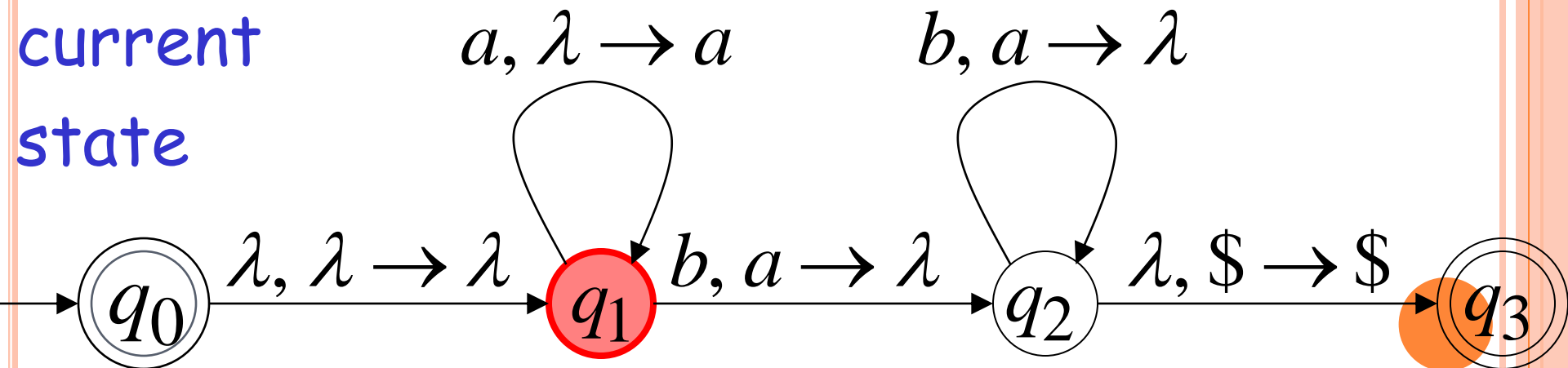
Rejection Example: Time 1

Input



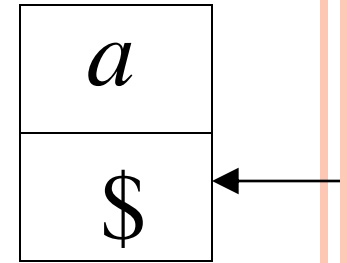
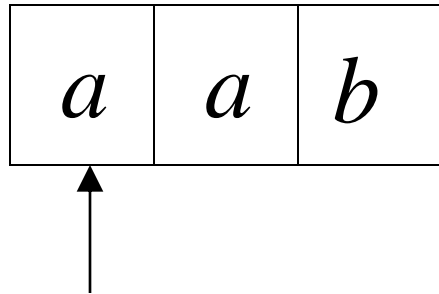
Stack

current
state



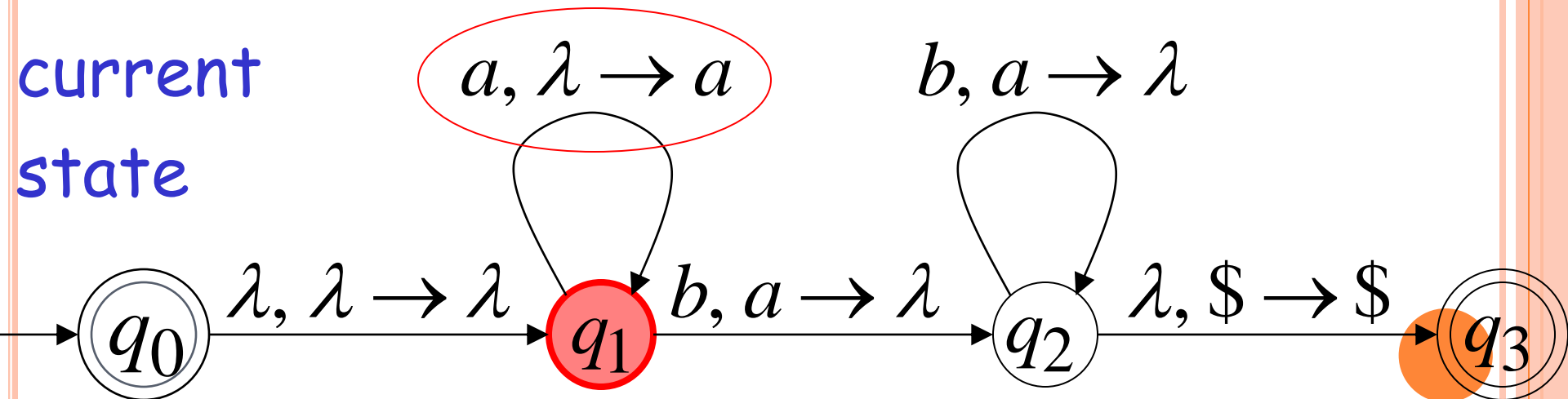
Rejection Example: Time 2

Input



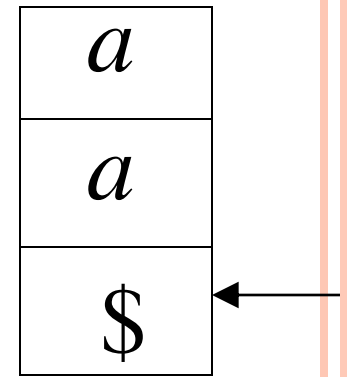
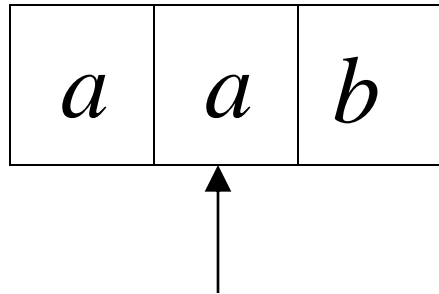
Stack

current
state



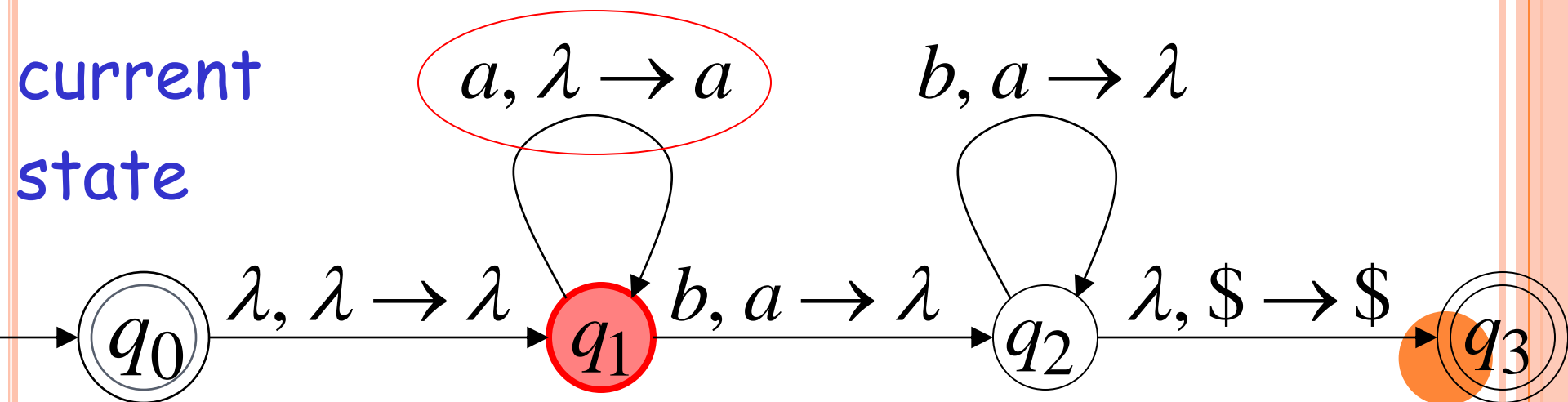
Rejection Example: Time 3

Input



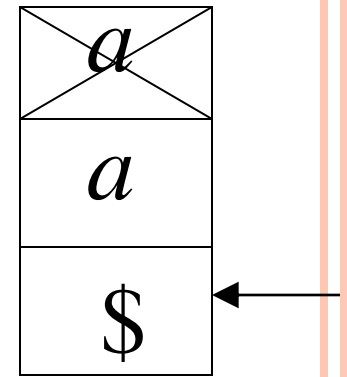
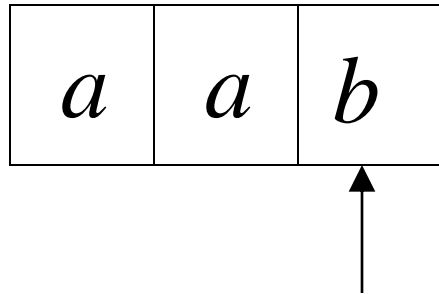
Stack

current
state



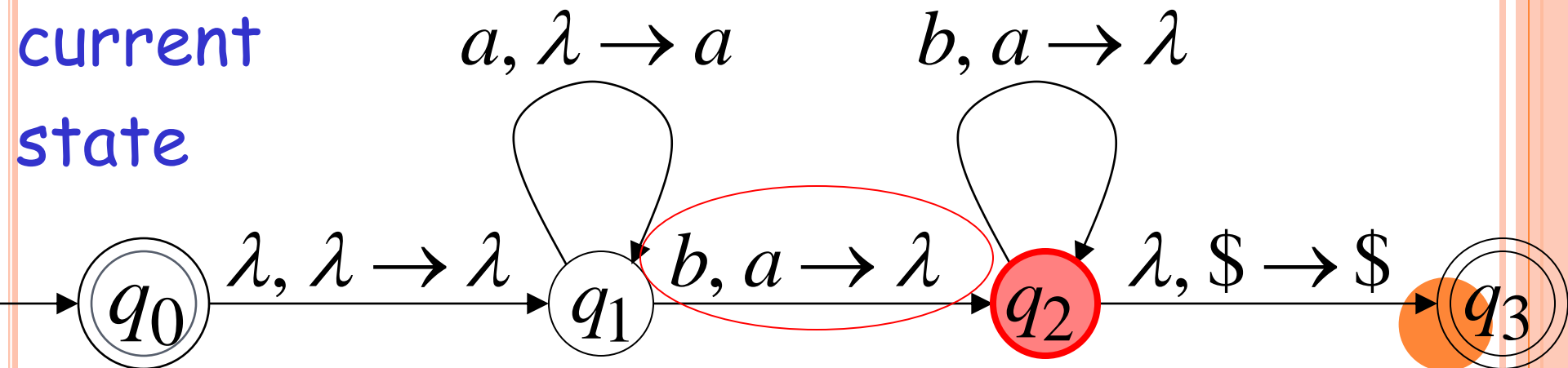
Rejection Example: Time 4

Input



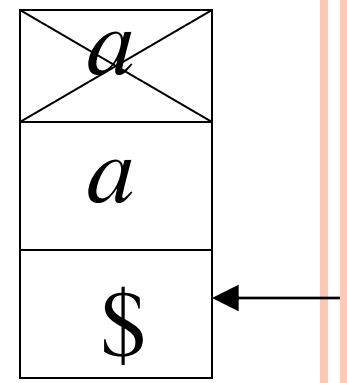
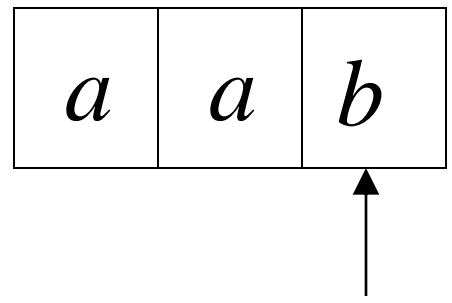
Stack

current
state



Rejection Example: Time 4

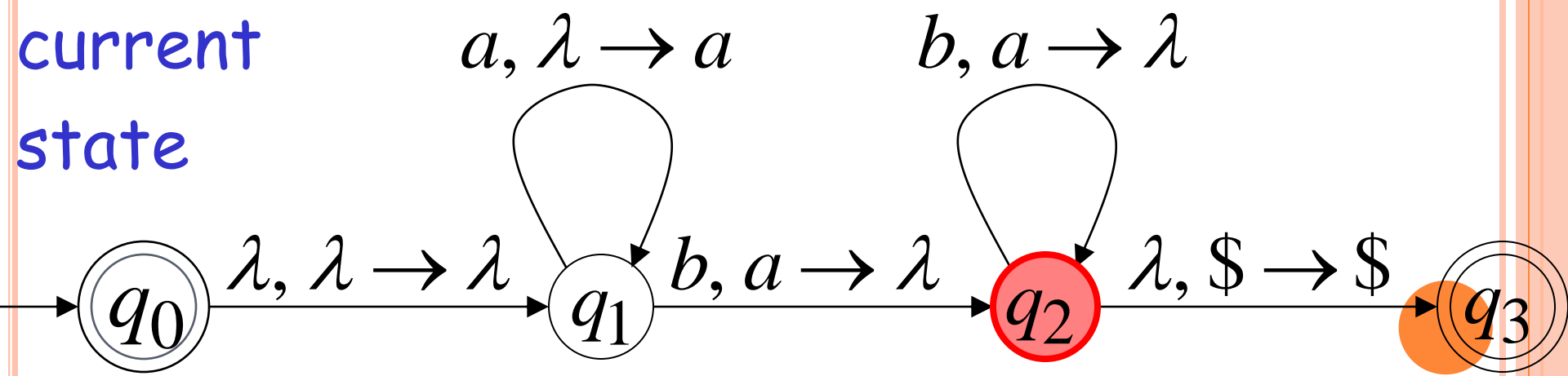
Input



Stack

reject

current state



There is no accepting computation for aab

The string aab is rejected by the PDA

